

SoK: Signatures With Randomizable Keys

*Sofía Celi, Scott Griffy, Lucjan Hanzlik,
Octavio Perez Kempner, Daniel Slamanig*

Brave Software, Brown University, CISPA,
NTT Social Informatics Laboratories, AIT Austrian Institute of Technology

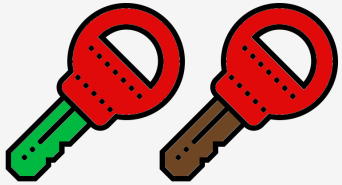


Digital Signatures





Digital Signatures



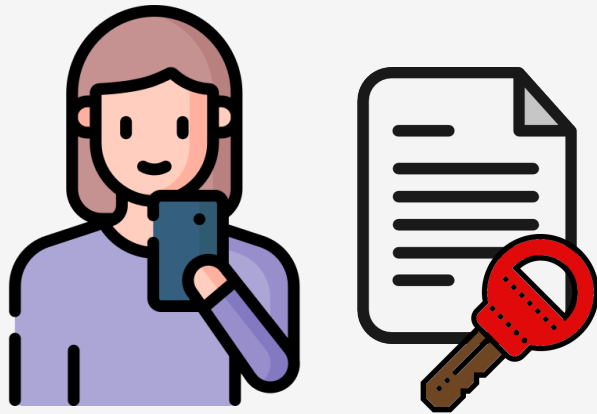


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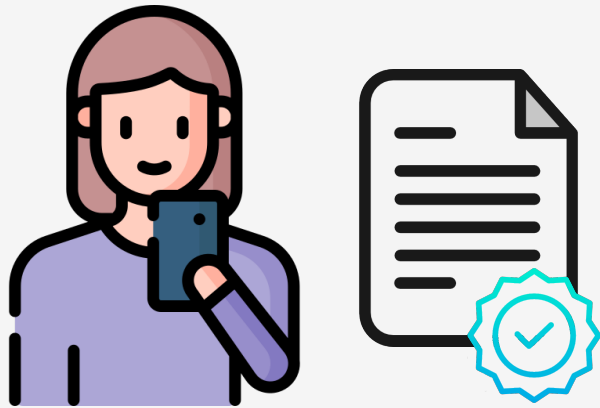


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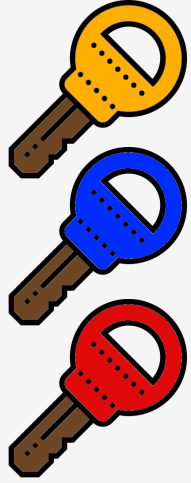


WebAuthn Application



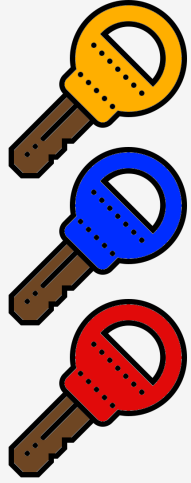


WebAuthn Application





WebAuthn Application



Can we do better?





Digital Signatures: Single Master Key



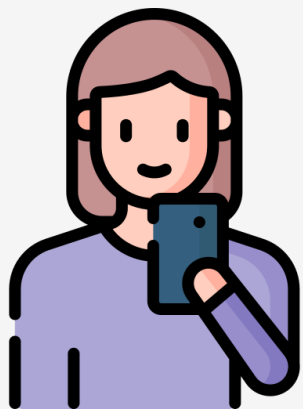


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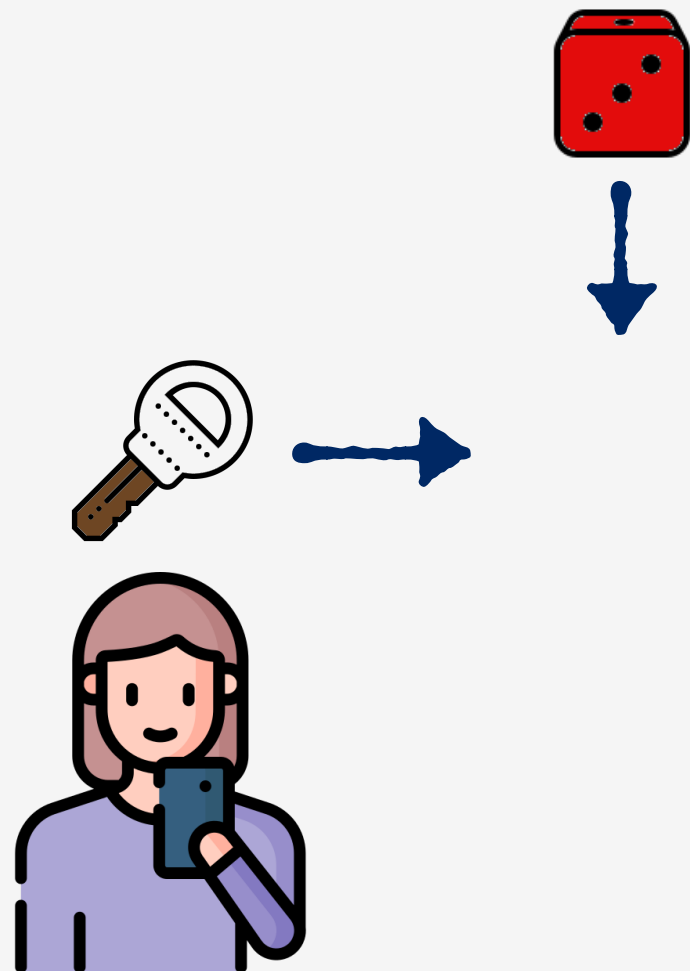


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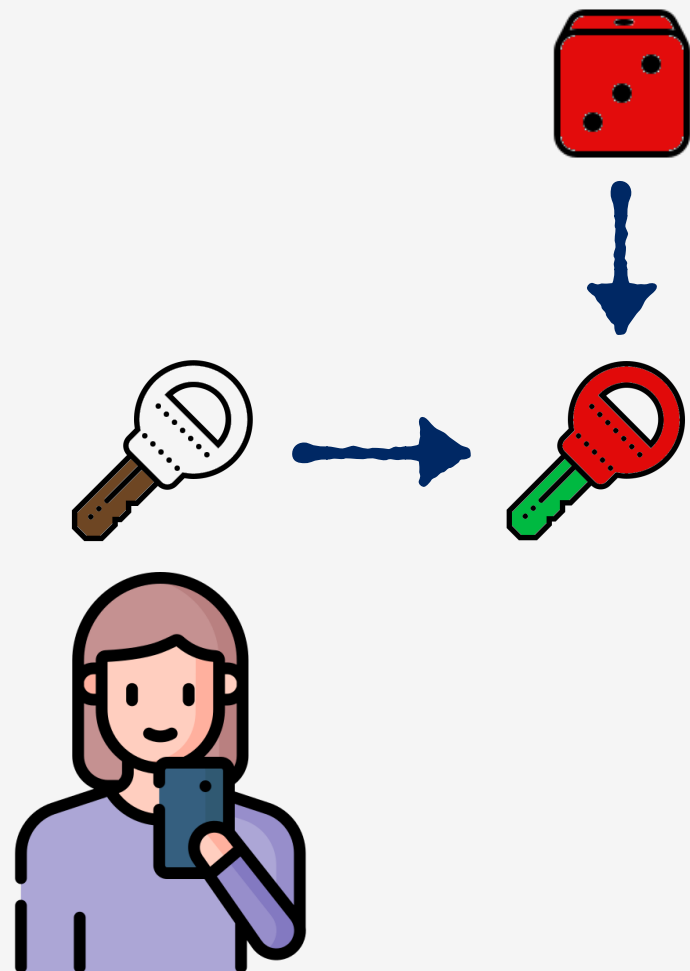


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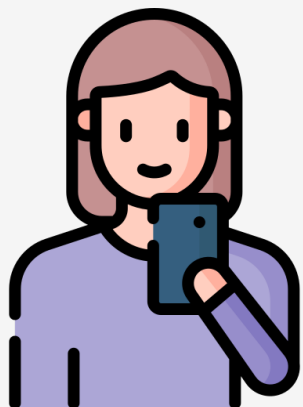


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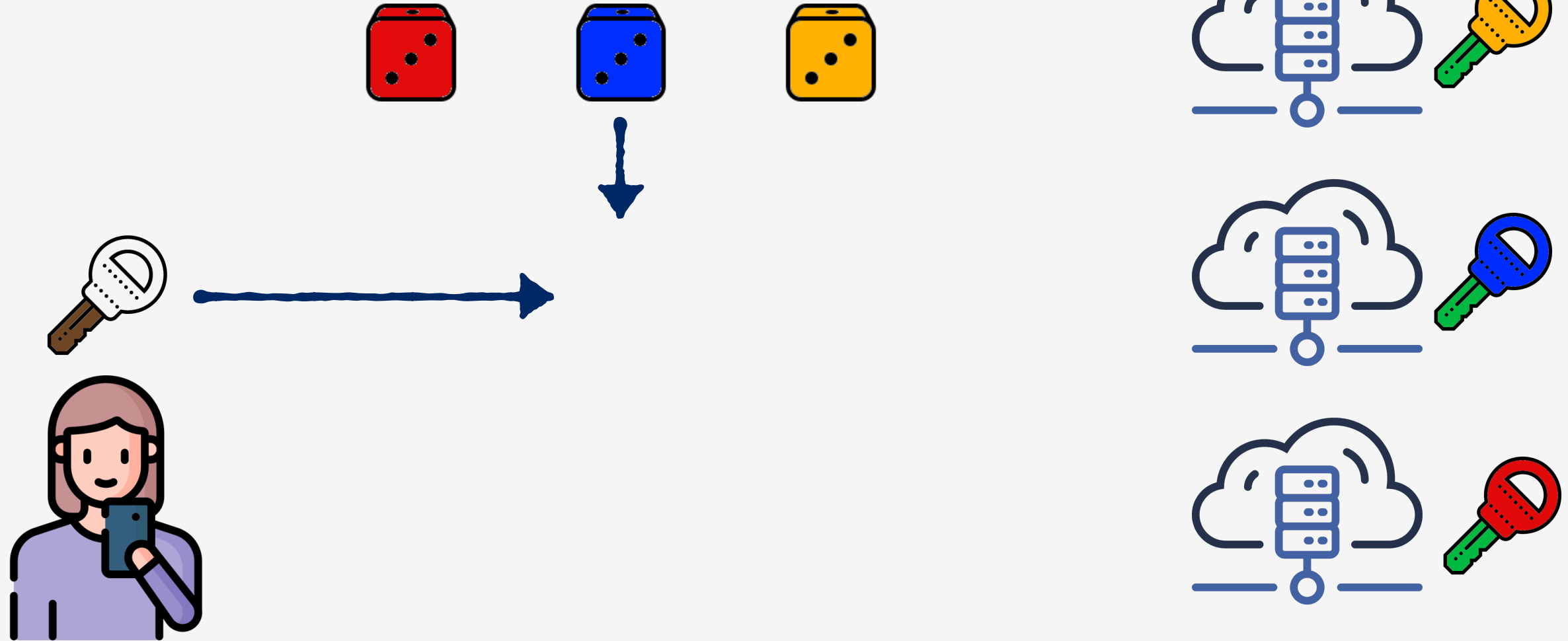


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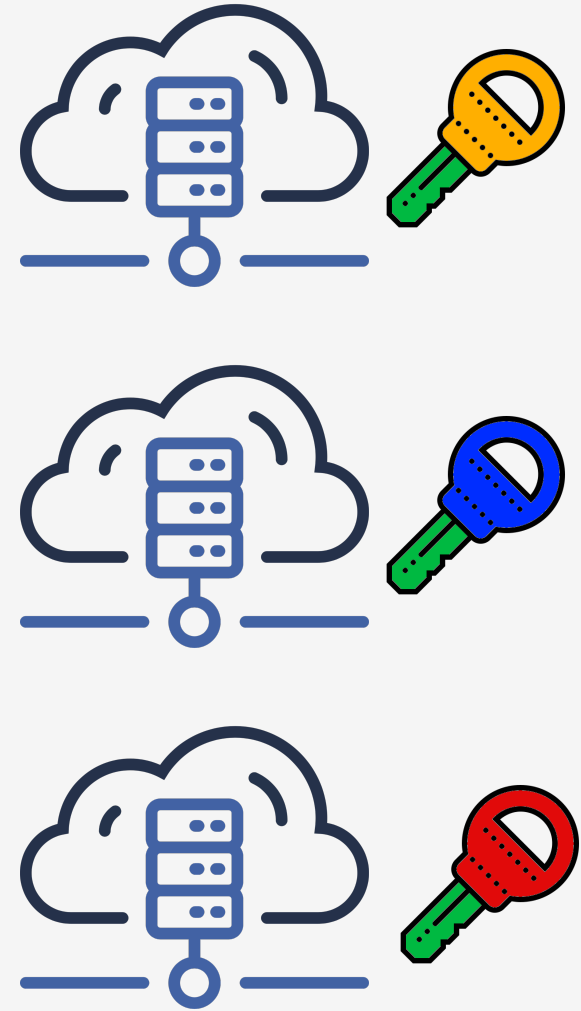
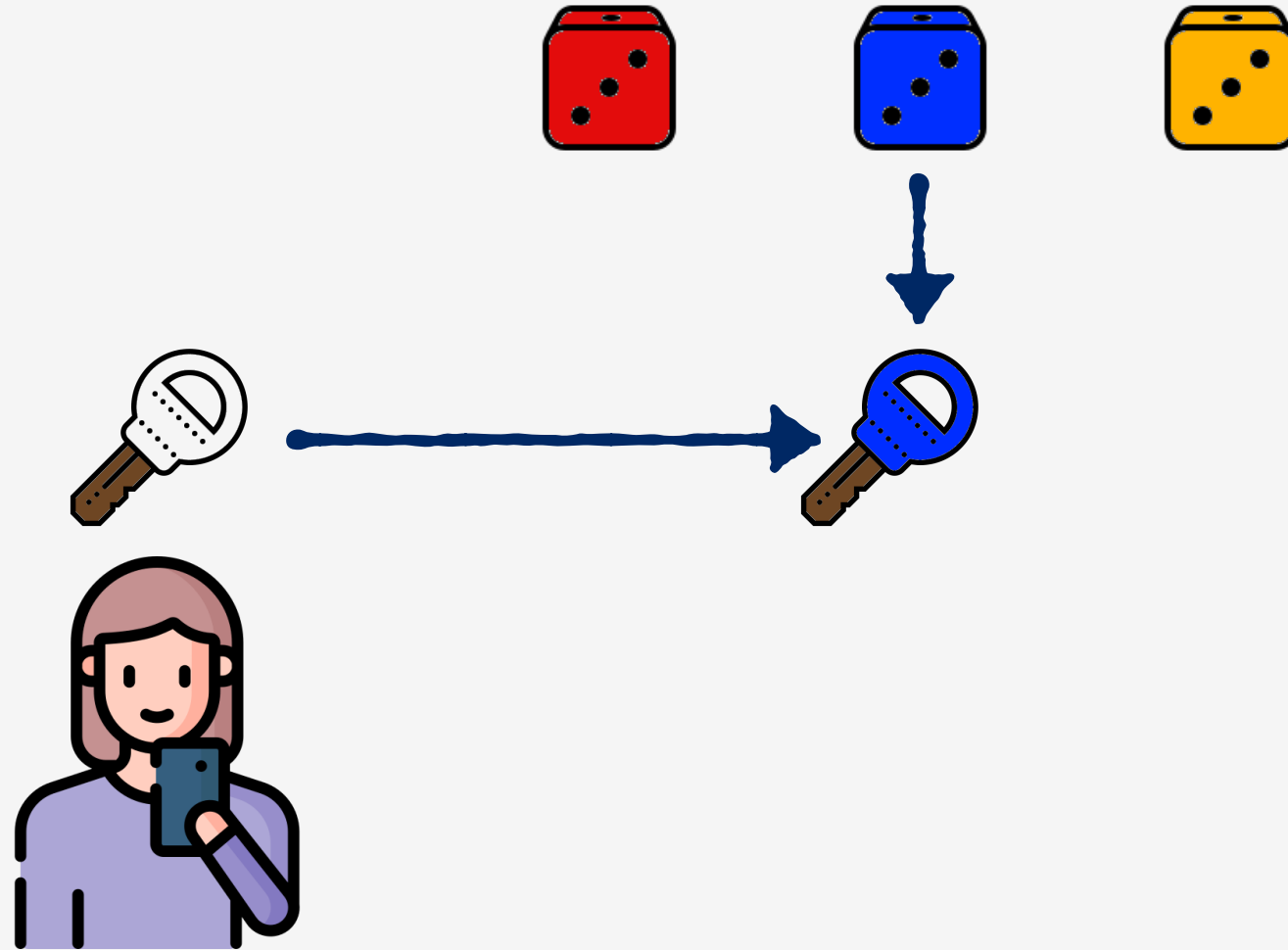


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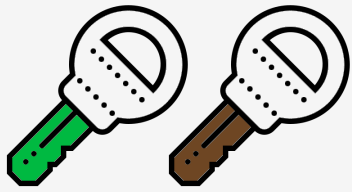




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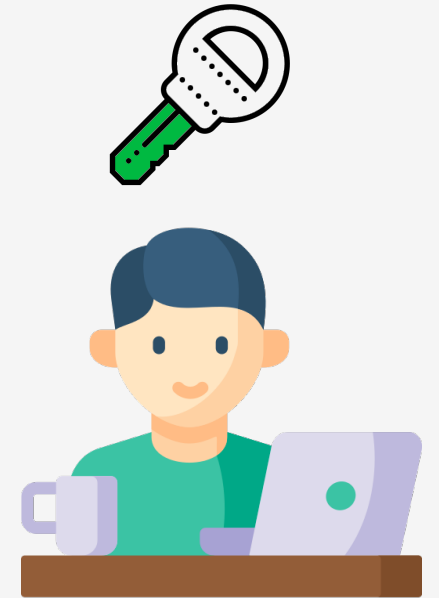
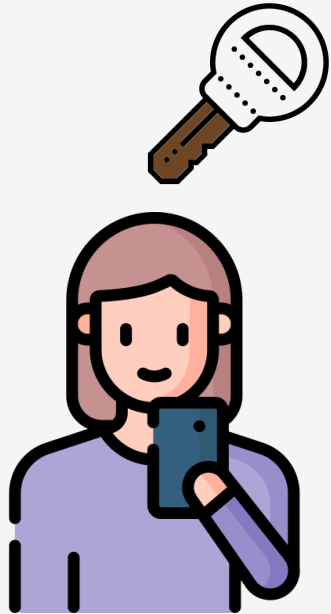


Digital Signatures: Public Randomization



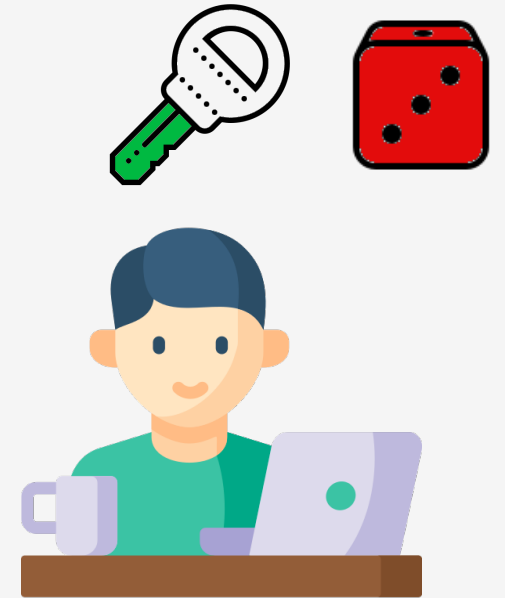
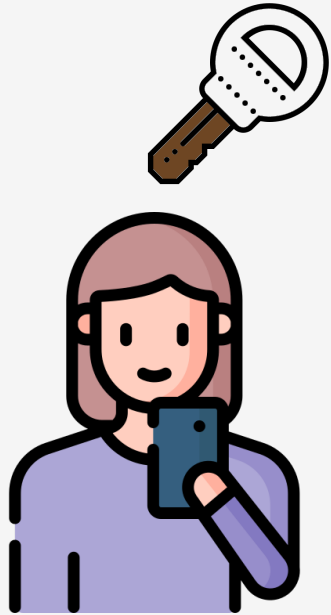


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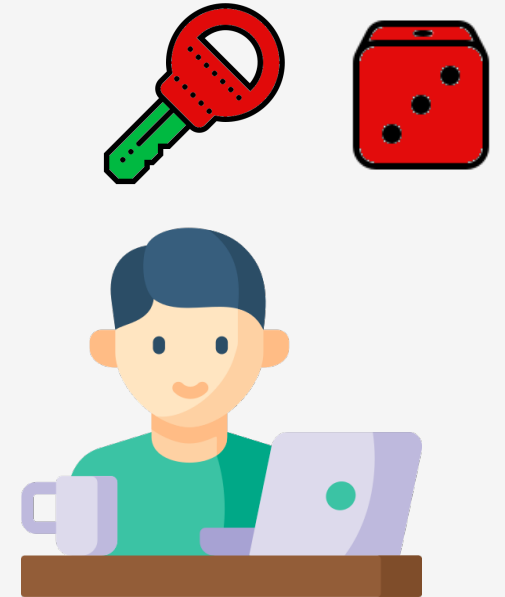
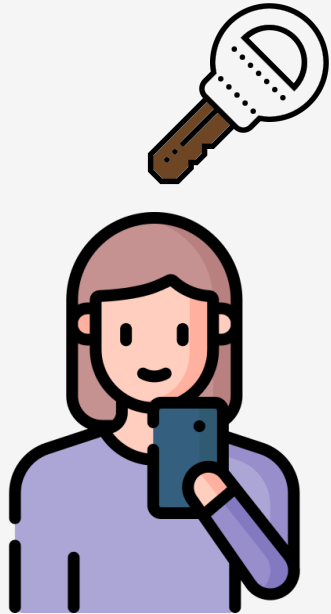


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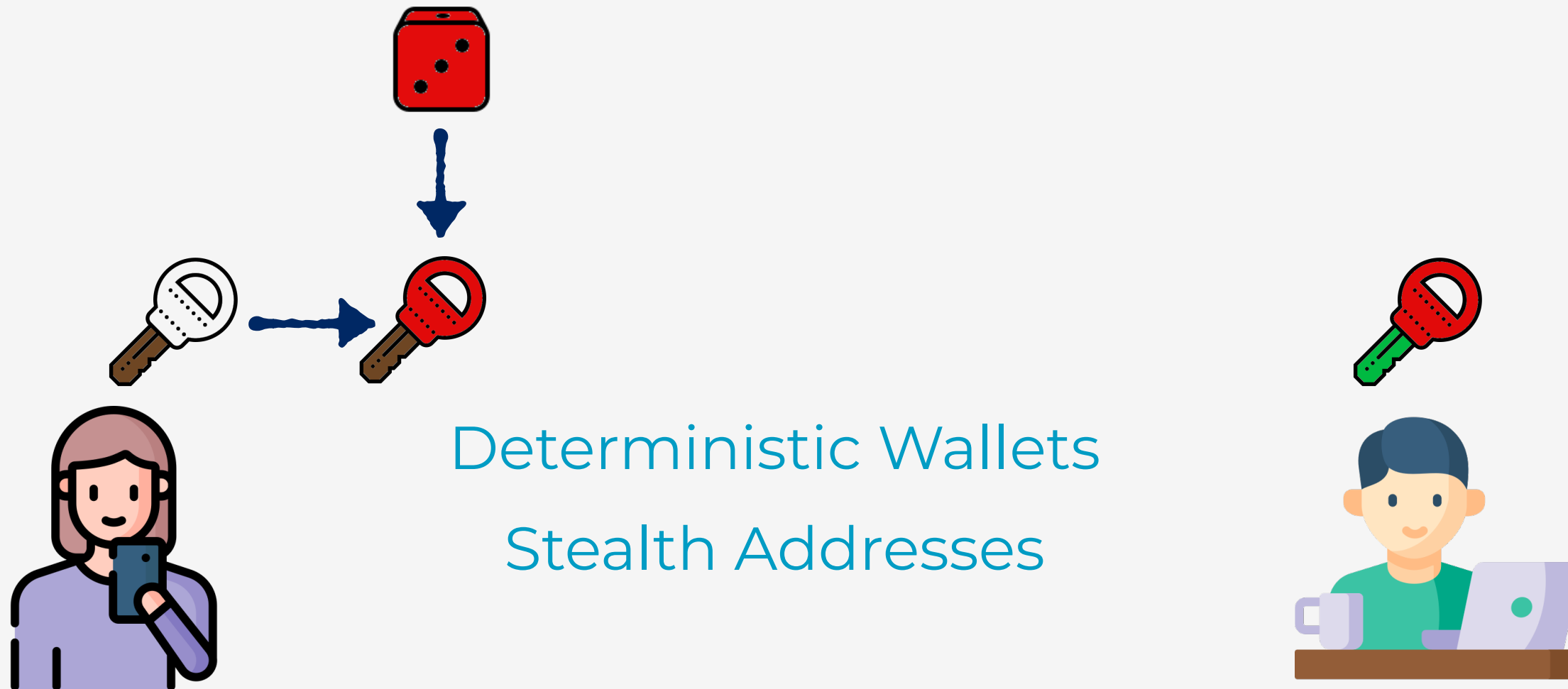


Digital Signatures: Public Randomization



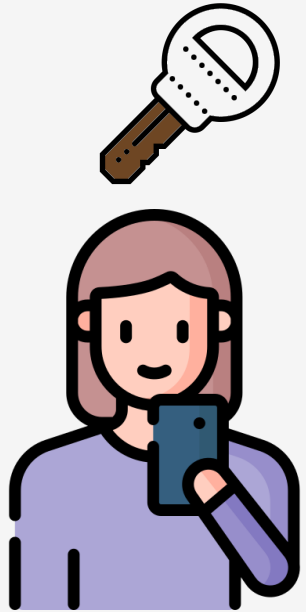


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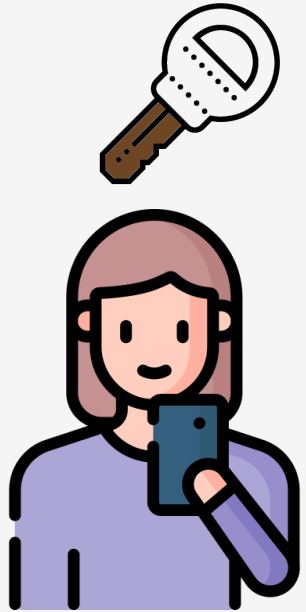


Digital Signatures: Randomizing Signature





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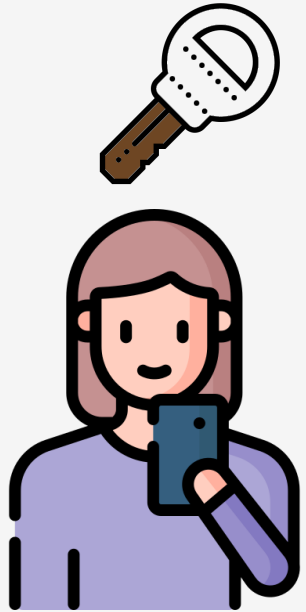


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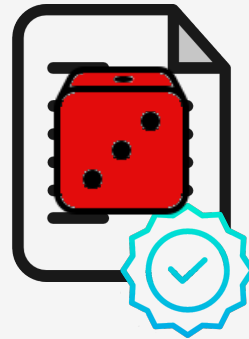
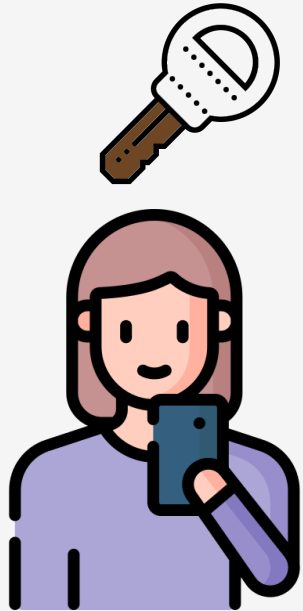


Digital Signatures: Randomizing Signature



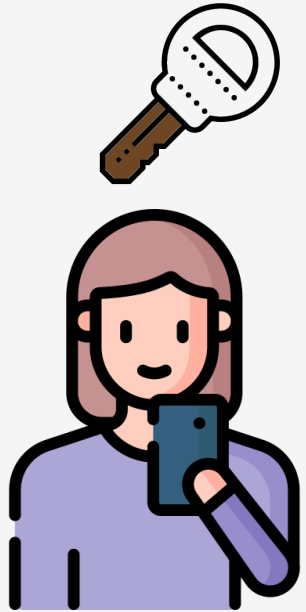


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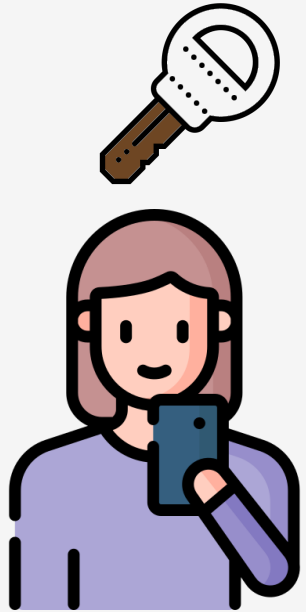


Digital Signatures: Randomizing Signature





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Digital Signatures: Different Notions

mercurial key-homomorphic
key-blinding re-randomizable flexible-public keys
signatures
equivalence-classes
homomorphic randomizable



Digital Signatures: Different Notions



* Sok (in Polish) = Juice



Our Contribution

- Introduce signatures with randomizable keys
 - extend digital signatures
 - parametrizable security properties
- Revisit prior work and how it relates to our syntax and model
- Show what is required for specific applications



New Algorithms: Randomization and Adaptation



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- Separate algorithms to randomize secret/public key (RandSK & RandPK)
 - take as input the original key and key randomizer
 - outputs randomized key, e.g., $pk' = T(pk, r)$



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- Separate algorithms to randomize secret/public key (RandSK & RandPK)
 - take as input the original key and key randomizer
 - outputs randomized key, e.g., $pk' = T(pk, r)$
- Optional adaptation algorithm
 - takes as input signature, public key and key randomizer
 - outputs signature valid under pk'
 - adapted signatures look like fresh signatures - (perfect) adaptation



New Security Properties

- Unforgeability - no forged signatures
- Unlinkability - randomized public keys are not linkable to original ones
- Unextractability - cannot go back to original public key even knowing the key randomizer



Security Properties: α -Unforgeability

(sk, pk)



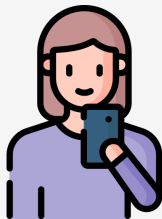
pk





Security Properties: α -Unforgeability

(sk, pk)



get key randomizer



r

pk



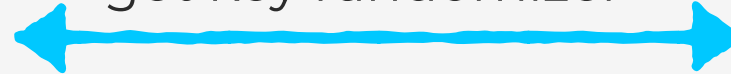


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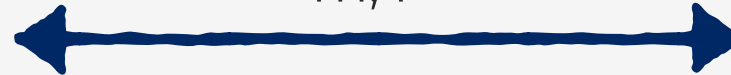


r

pk



m, r



σ valid under $pk' = T(pk, r)$

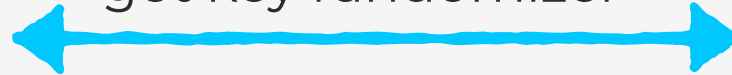


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from Oracle (if $\alpha=0$)



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⋮

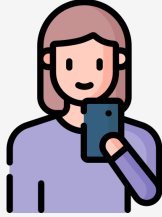
m^*, σ^*, r^*





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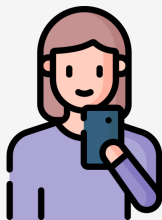


Adv wins iff

- σ^* valid for $pk^* = T(pk, r^*)$
- r^* from Oracle (if $\alpha=0$)



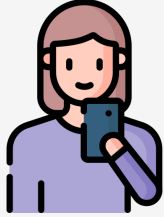
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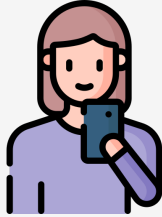
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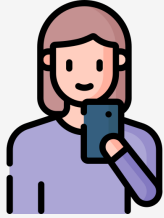
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picks bit b and r

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Security Properties: (α, β, γ) -Unlinkability

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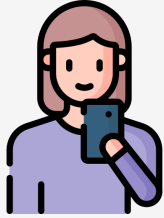
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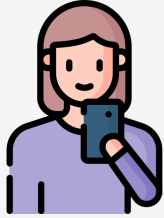
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signing oracle



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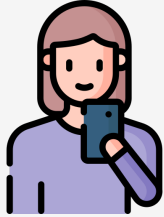
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Adv wins iff $b = b^*$



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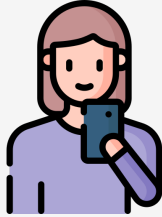
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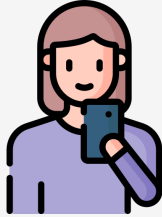
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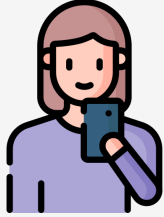




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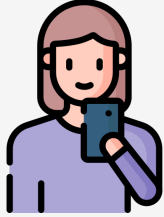




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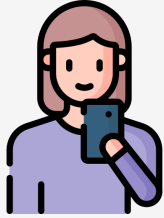
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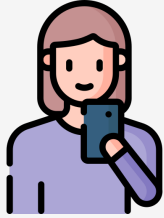
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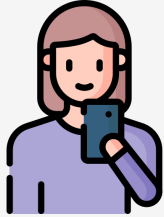
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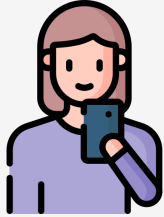


access to all keys



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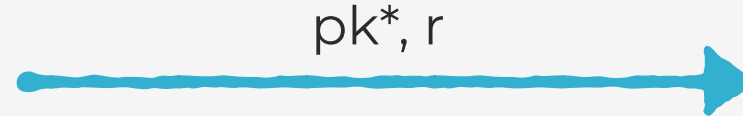
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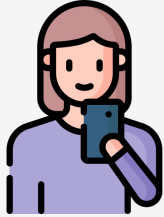


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BLS Signatures - Next Talk

- Support perfect adaptation
- Unforgeable against malicious randomizer (1-UNF)
- (1,1,3)-Unlinkable but are not unextractable
- Can be used for Deterministic Wallets and Stealth Addresses which can work with (0,0,3)-UNL and 1-UNF



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See paper for full systematization



Thank you for your attention

